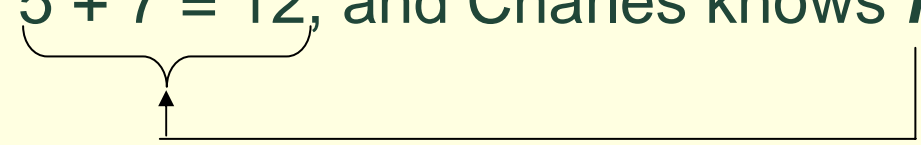


# Semantic pre-processing of anaphoric references

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# Using TIL (of course)

(S) “5 + 7 = 12, and Charles knows *it*.”



- But Charles doesn't know T. He knows that the **procedure** of calculating  $5+7=12$  yields T.
- Thus the meaning of (S) is a two-phase instruction that comes down to this:
  - i. **Pre-processing** (of the meaning of the embedded clause): execute *the substitution based on the meaning of the antecedent (here:  $5+7=12$ ) for the anaphoric variable (here: *it*)*;
  - ii. **Execute the result (a propositional construction) again to obtain a proposition.**

# *Method of analysis*

(S) “5 + 7 = 12, and Charles knows *it*.”

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a) Type-theoretical analysis.

**5,7,12/τ**; τ is the type of real numbers.

**+/(τττ)**; the function of adding that maps number-pairs (ττ) to a number (of type τ).

**=/(οττ)**; the relation of identity on numbers.

**∧/(οοο)**; conjunction: mapping couples of truth values (type (οο)) to truth values (ο).

**Charles/ι**; ι is the type of individuals.

**Know/(οι\*<sub>1</sub>)<sub>τω</sub>**; the relation-in-intension (τω) of an individual (ι) to a construction (procedure of type \*<sub>1</sub>).

**it/\*<sub>2</sub> → \*<sub>1</sub>**; the anaphoric variable.

**Subl(\*<sub>2</sub>\*<sub>2</sub>\*<sub>2</sub>\*<sub>2</sub>)**—the substitution function: applied on constructions C<sub>1</sub>, C<sub>2</sub>, C<sub>3</sub>, it yields a construction C<sub>4</sub> that is the result of substituting C<sub>1</sub> for C<sub>2</sub> into C<sub>3</sub>.

# *Method of analysis*

(S) “5 + 7 = 12, and Charles knows *it*.”

b) Synthesis.

‘5+7=12’ →

$[^0 = [^0 + ^0 5 ^0 7] ^0 12];$

‘Charles knows it’ →

$\lambda w \lambda t [^0 Know_{wt} ^0 Charles it];$

Gloss: **Constructions** are instructions (**procedures**);

*Atomic constructions* (supply entities that the composed constructions operate on):

*Trivialisation* ( $^0 =, ^0 +, ^0 5, ^0 7, ^0 12$ ), *Variables* ( $x, p, it, he, \dots$ );

*Composed constructions* consist of constituents:

*Composition* (like  $[^0 + x ^0 1]$ ), *Closure* ( $\lambda x [^0 + x ^0 1]$ ), *Double Execution*  $^2 C$ ;

(S) →  $\lambda w \lambda t [[^0 = [^0 + ^0 5 ^0 7] ^0 12] \wedge$

$^2 [^0 Sub ^0 [^0 = [^0 + ^0 5 ^0 7] ^0 12] ^0 it$

$^0 [\lambda w \lambda t [^0 Know_{wt} ^0 Charles it]]_{wt}]$

# *Method of analysis*

a) *Pre-processing (1<sup>st</sup> execution):*

*Sub(stitutes) the meaning of the antecedent  
( ${}^0[{}^0=[{}^0+ {}^05 {}^07] {}^012]$ ) for the anaphoric variable  
(it) into the meaning of the embedded clause  
( $[\lambda w \lambda t [{}^0Know_{wt} {}^0Charles\ it]$ );*

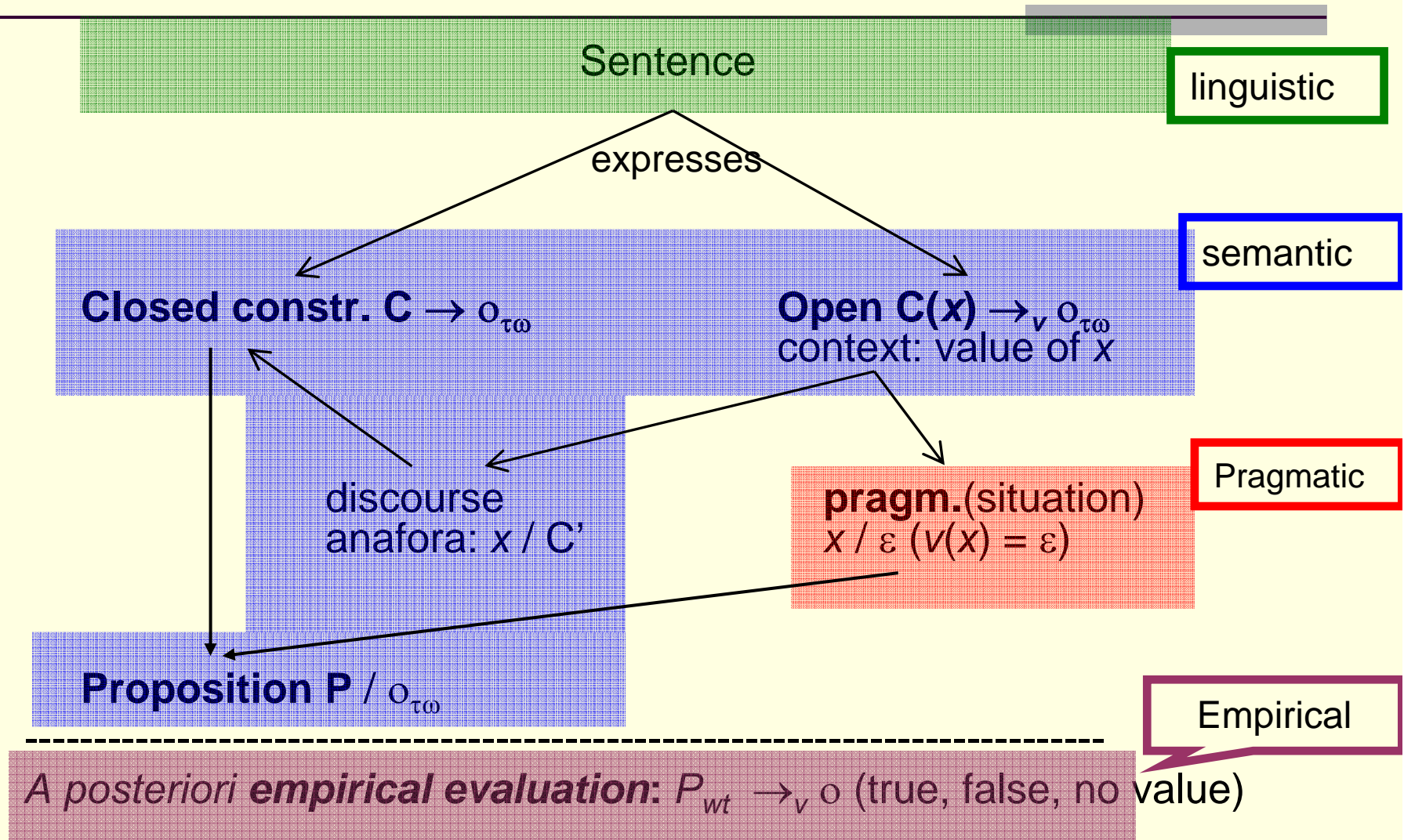
b) *2<sup>nd</sup> Execution:*

*of the adjusted (now closed) embedded  
clause constructs the proposition (that  
Charles knows ...)*

c) *Intensional descent:  $[[C\ w]\ t]$ ,  $C_{wt}$  yields the  
truth value*

# Semantic conception of TIL

Levels:



# *de re* attitudes & anaphora

“Charles is looking for the Mayor of Dunedin”.

- understood as uttered in a situation where Charles knows who the Mayor is, and is striving to *locate* this individual:

“Charles is looking for the Mayor of Dunedin, namely the location of *him*”.

- $\lambda w \lambda t [{}^0 \text{Look\_for}_{wt} {}^0 \text{Charles}$   
 ${}^2 [{}^0 \text{Sub} [{}^0 \text{Tr} [{}^0 \text{Mayor\_of}_{wt} {}^0 D]] {}^0 \text{him} {}^0 [\lambda w \lambda t [{}^0 \text{Loc}_{wt} \text{him}]]]$

v-constructs Trivialisation of the individual (if any) who occupies the Mayor office.

Types:  $\text{Look\_for}/(\circ \iota \mu_{\tau\omega})_{\tau\omega}$ ;  $\text{Tr}/(*_1 \iota)$ ;  $\text{Charles}/\iota$ ;  $\text{Mayor\_of}/(\iota \iota)_{\tau\omega}$ ;  $D(\text{unedin})/\iota$ ;  $\text{him}/*_1 \rightarrow \iota$ ;  $\text{Loc}/(\mu \iota)_{\tau\omega}$ .

# Donkey sentences

- “If somebody has got a new car then *he* often washes *it*.”
- Embedded clause:  $\lambda w \lambda t [{}^0 \text{Wash}_{wt} \text{ he it}]$ 
  - $\lambda w \lambda t [{}^0 \text{Freq}_t \lambda t' [{}^0 \text{Wash}_{wt'} \text{ he it}]]$
- *Problem*: how to understand the sentence?
- If somebody owns more than one new car, does he wash *all* or *some* of them?



# Donkey sentences

Peter **Geach** (*NC* new car):

$\forall x \forall y ((NC(y) \wedge Has(x, y)) \rightarrow Wash(x, y)).$

Bertrand **Russell** ('a new car' is an *indefinite description*):

$\forall x (\exists y (NC(y) \wedge Has(x, y)) \rightarrow Wash(x, y)).$

But the last occurrence of the variable  $y$  (marked in red) is free in this formula—out of the scope of the existential quantifier supposed to bind it.

**Neale** (*restricted quantifiers*):

[every  $x$ : man  $x$  and [a  $y$ : new-car  $y$ ]( $x$  owns  $y$ )]([whe  $z$ : car  $z$  and  $x$  owns  $z$ ] ( $x$  often washes  $z$ )).

Neale takes into account that the sentence is true even if a man owns *more than one* new car: his abbreviation 'whe  $F$ ' stands for 'the  $F$  or the  $F$ s'.

# Donkey sentences

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(D1) “**Anybody who owns some new cars often washes *all of them*** [each of the new cars he owns].”

(D2) “**Anybody who owns some new cars often washes *some of them*** [some of the new cars he owns].”

■ *Types*:  $Own/(o11)_{\tau\omega}$ ;  $Wash/(o11)_{\tau\omega}$ ;  $NC$  (being a new car)/ $(o1)_{\tau\omega}$ ;  $x, y, he, it \rightarrow 1$ .

■ We need another type of quantifiers: *Some, All*

# Some, All quantifiers

- *Some, All* of type  $((o(o1))(o1))$ .
- *Some* is a function that associates the argument—a set  $S$ —with the set of all those sets which have a non-empty intersection with  $S$ .
- *All* is a function that associates the argument—a set  $S$ —with the set of all those sets which contain  $S$  as a subset.
- $\lambda w \lambda t [[{}^0\text{Some } {}^0\text{Student}_{wt}] {}^0\text{Happy}_{wt}]$ .

# Donkey sentences

(D1'')

$$\lambda w \lambda t \left[ {}^0 \forall \lambda x \left[ \left[ {}^0 \text{Man}_{wt} x \right] \wedge \right. \right. \\ \left. \left[ {}^0 \exists \lambda y \left[ \left[ {}^0 \text{NC}_{wt} y \right] \wedge \left[ {}^0 \text{Own}_{wt} x y \right] \right] \right] \supset \right. \\ \left. {}^2 \left[ {}^0 \text{Sub } {}^0 \left[ \lambda y \left[ \left[ {}^0 \text{NC}_{wt} y \right] \wedge \left[ {}^0 \text{Own}_{wt} x y \right] \right] \right] {}^0 \text{them} \right. \right. \\ \left. \left[ {}^0 \text{Sub } {}^0 x {}^0 \text{he} \right. \right. \\ \left. \left. {}^0 \left[ \lambda w' \lambda t' \left[ \left[ {}^0 \text{All them} \right] \lambda it \left[ {}^0 \text{Wash}_{w't'} \text{he it} \right] \right] \right] \right]_{wt} \right].$$

- Gloss: “For every man, if the man owns some new cars then **all of them** [i.e., **the new cars owned**] are washed by him [the man x].”

# Donkey sentences

(D2'')

$$\lambda w \lambda t \left[ {}^0 \forall \lambda x \left[ \left[ {}^0 \text{Man}_{wt} x \right] \wedge \right. \right. \\ \left. \left[ {}^0 \exists \lambda y \left[ \left[ {}^0 \text{NC}_{wt} y \right] \wedge \left[ {}^0 \text{Own}_{wt} x y \right] \right] \right] \supset \right. \\ \left. {}^2 \left[ {}^0 \text{Sub } {}^0 \left[ \lambda y \left[ \left[ {}^0 \text{NC}_{wt} y \right] \wedge \left[ {}^0 \text{Own}_{wt} x y \right] \right] \right] {}^0 \text{them} \right. \right. \\ \left. \left[ {}^0 \text{Sub } {}^0 x {}^0 \text{he} \right. \right. \\ \left. \left. {}^0 \left[ \lambda w' \lambda t' \left[ \left[ {}^0 \text{Some them} \right] \lambda it \left[ {}^0 \text{Wash}_{w't'} \text{he it} \right] \right] \right] \right]_{wt} \right].$$

- Gloss: “For every man, if the man owns some new cars then **some of them** [i.e., **the new cars owned**] are washed by him [the man x].”

# Compositional analysis

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Gabriel Sandu formulates (1997):

- there is a one-to-one mapping of the surface structure of a sentence of (a fragment of) English into its logical form which preserves the left-to-right ordering of the logical constants
- the mapping preserves the nature of the lexical properties of the logical constants, in the sense that ***an indefinite is translated by an existential quantifier***, etc.

# Compositional analysis (of D)

In the interest of disambiguation, we actually analysed two variants of the original sentence (a man vs. every man).

The analysis of the antecedent clause “**A man has a new car**” should be as follows:

(NC)  $\lambda w \lambda t [{}^0 \exists \lambda x y [ [{}^0 \text{Man}_{wt} x] \wedge [{}^0 \text{NC}_{wt} y] \wedge [{}^0 \text{Own}_{wt} x y] ]]$ .

■ Additional type:  $\exists / (o(o\iota))$ .

The consequent of (D) expresses that **all the couples <he, it> are such that he Washes it**. Using a variable *couples*/\* $_1 \rightarrow (o\iota)$ , we have:

(AC)  $\lambda w \lambda t [ [{}^0 \text{All couples}] \lambda he it [{}^0 \text{Wash}_{wt} he it] ]]$ .

# Compositional analysis (of D)

- Composing (NC) with (AC), we substitute the set of couples ... for the variable *couples*:

$$(D') \quad \lambda w \lambda t \left[ \left[ \exists \lambda xy \left[ \left[ {}^0 \text{Man}_{wt} x \right] \wedge \left[ {}^0 \text{NC}_{wt} y \right] \wedge \left[ {}^0 \text{Own}_{wt} x y \right] \right] \supset \right. \right. \\ \left. \left. {}^2 \left[ {}^0 \text{Sub } {}^0 \left[ \lambda xy \left[ \left[ {}^0 \text{Man}_{wt} x \right] \wedge \left[ {}^0 \text{NC}_{wt} y \right] \wedge \left[ {}^0 \text{Own}_{wt} x y \right] \right] \right] \right] {}^0 \text{couples} \right. \right. \\ \left. \left. {}^0 \left[ \lambda w \lambda t \left[ \left[ {}^0 \text{All couples} \right] \lambda he it \left[ {}^0 \text{Wash}_{wt} he it \right] \right] \right] \right]_{wt} \right].$$

The Dynamic Predicate Logic (DPL): passing on binding.

(D') – the semantics of this mechanism.

*Variables he and it* are bound in (D'), but not directly bound by the existential quantifier.

Technically, they are bound by Trivialization; semantically, they are bound by the condition that the pairs of individuals they v-construct have to belong to the set mentioned by the antecedent clause.



# Discourse Representation Theory (Kamp & Reyle)

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- dynamic interpretation of natural language, where each sentence is interpreted within a certain discourse, which is a sequence of sentences uttered by the same speaker.
- the problem of anaphoric links crossing the sentence boundary.
- *Pressing question*: to determine the respective antecedent to which the anaphoric pronoun refers.
- first-order theory; only expressions denoting individuals introduce the so-called *discourse referents*, i.e., free variables that are updated when interpreting the discourse.

# TIL and Discourse Representation

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- TIL: higher-order, procedural
- not only individuals, but ***entities of any type***, like *properties, propositions, relations-in-intension*, and even *constructions* can be ***linked to anaphoric variables***.
- Moreover, strong typing makes it possible to ***determine*** the respective ***type-appropriate antecedent(s)***.

# Outline of Implementation method (first proposed by J. Křetínský)

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- For each type  $(\iota, (\sigma\iota)\tau\omega, \sigma\tau\omega, (\sigma\iota(\sigma\iota)\tau\omega)\tau\omega, (\sigma\iota\iota)\tau\omega, *n, \dots)$  we create the ***list of free discourse variables***. They are dynamic (program-like) variables.
- The method substitutes their content for anaphoric (logical) variables to complete the meaning of anaphoric clauses.
- Each closed constituent of a resulting construction becomes an updated value of the respective (type-appropriate) free discourse-referent variable.
- In this way the discourse variables are gradually updated.

# Example: dialogue between *Adam*, *Berta* and *Cecil*.

- *Adam to Cecil*: “Berta is coming. **She** is looking for a parking”.

‘Inform’ message content:

$\lambda w \lambda t \text{ [[}^0\text{Coming}_{wt} \text{ }^0\text{Berta}].$

- (Relevant) discourse variables updates:

$\text{ind} := ^0\text{Berta}; \text{pred} := ^0\text{Coming}; \text{prop} := \lambda w \lambda t \text{ [[}^0\text{Coming}_{wt} \text{ }^0\text{Berta}];$

$\lambda w \lambda t \text{ }^2\text{[}^0\text{Sub } \text{ind} \text{ }^0\text{she } ^0\text{[}^0\text{Looking\_for}_{wt} \text{ she } ^0\text{Parking}]} \Rightarrow$  (is transformed into)

$\lambda w \lambda t \text{ [}^0\text{Looking\_for}_{wt} \text{ }^0\text{Berta } ^0\text{Parking}].$

- (Relevant) discourse variables updates:

$\text{rel1} := ^0\text{Looking\_for}; \text{pred} := ^0\text{Parking};$

$\text{prop} := \lambda w \lambda t \text{ [}^0\text{Looking\_for}_{wt} \text{ }^0\text{Berta } ^0\text{Parking}];$

$\text{prof} := \lambda w \lambda t \lambda x \text{ [}^0\text{Looking\_for}_{wt} \text{ } x \text{ } ^0\text{Parking}];$

# Example: dialogue between *Adam*, *Berta* and *Cecil*.

- *Cecil to Adam*: “**So** am I.”

$\lambda w \lambda t \ ^2 [^0 \text{Sub } \mathbf{prof} \ ^0 \mathbf{so} \ ^0 [so_{wt} \ ^0 \text{Cecil}]] \Rightarrow$   
 $\lambda w \lambda t \ [^0 \text{Looking\_for}_{wt} \ ^0 \text{Cecil} \ ^0 \text{Parking}].$

- (Relevant) discourse variables updates:

$ind := ^0 \text{Cecil}; rel_1 := ^0 \text{Looking\_for}; pred := ^0 \text{Parking};$

- *Adam to both*: “There is a free parking at  $p_1$ ”.

$\lambda w \lambda t \ \exists x \ [ [ [^0 \text{Free } ^0 \text{Parking}]_{wt} \ x] \wedge [^0 \text{At}_{wt} \ x \ ^0 p_1] ]$

- (Relevant) discourse variables updates:

$loc := ^0 p_1; pred := [^0 \text{Free } ^0 \text{Parking}];$

$prop := \lambda w \lambda t \ \exists x \ [ [ [^0 \text{Free } ^0 \text{Parking}]_{wt} \ x] \wedge [^0 \text{At}_{wt} \ x \ ^0 p_1] ].$

# Example: dialogue between *Adam*, *Berta* and *Cecil*.

- *Berta to Adam*: “What do you mean by free parking?”

‘**Query**’ message content:

$$\lambda w \lambda t [{}^0\text{Refine}_{wt} \text{ } {}^0[\text{}^0\text{Free } \text{}^0\text{Parking}]]$$

- (Relevant) discourse variables updates:

$$\text{constr.} = \text{}^0[\text{}^0\text{Free } \text{}^0\text{Parking}]$$

- *Adam to Berta*: “Free parking is a parking and some parts of it are not occupied”.

‘**Reply**’ message content:

$$[{}^0\text{Free } \text{}^0\text{Parking}] =$$

$$[\lambda w \lambda t \lambda x [{}^0\text{Parking}_{wt} x] \wedge$$

$$\exists y [{}^0\text{Part\_of}_{wt} y x] \wedge \neg [{}^0\text{Occupied}_{wt} y]]]$$

# Example: dialogue between *Adam*, *Berta* and *Cecil*.

- *Berta to Adam*: “I don’t believe *it*. I have just been *there*”.

‘Inform’ message content (first sentence):

$$\lambda w \lambda t \ [^2[^0\text{Sub } \mathbf{prop} \ ^0\mathbf{it} \ ^0[\neg[^0\text{Believe}_{wt} \ ^0\text{Berta } \mathbf{it}]]]] \Rightarrow$$
$$\lambda w \lambda t \ \neg[^0\text{Believe}_{wt} \ ^0\text{Berta} \ [\lambda w \lambda t \ \exists x \ [ [[[^0\text{Free } \ ^0\text{Parking}]_{wt} \ x] \wedge \ [^0\text{At}_{wt} \ x \ ^0p_1]]]].$$

‘Inform’ message content (second sentence):

$$\lambda w \lambda t \ \exists t' \ [[t' \leq t] \wedge$$
$$\ ^2[^0\text{Sub } \mathbf{loc} \ ^0\mathbf{there} \ ^0[^0\text{Been\_at}_{wt'} \ ^0\text{Berta } \mathbf{there}]]] \Rightarrow$$
$$\lambda w \lambda t \ \exists t' \ [[t' \leq t] \wedge \ [^0\text{Been\_at}_{wt'} \ ^0\text{Berta} \ ^0p_1]].$$

# Concluding remarks

- Due to the procedural semantics, our **agents can learn new concepts** by asking the other agents.
- In our example, after receiving Adam's reply Berta learns the refined meaning of the 'free parking' predicate, i.e., she updates her knowledge base by the respective composed construction.
- Moreover, though our approach is as fine-grained as the syntactic approach of standard FIPA languages like KIF, the **content of agent's knowledge is not a piece of syntax**, but its **meaning** (i.e., TIL construction).
- And since the respective construction is what synonymous expressions (even of different languages) have in common, the **agents (should) behave in the same way independently of the language** in which their knowledge and ontology is expressed.
- For instance, if we switch to Czech, the underlying constructions are **identical**:  ${}^0[{}^0Free\ {}^0Parking] = {}^0[{}^0Volné\ {}^0Parkoviště]$ .



# Concluding remarks

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- Of course, improvements of the above method are straightforward.
- For instance, in the example we were substituting the last type-appropriate entity that received mention;
- If we wanted to take into account ambiguities of anaphoric references, we might store into the discourse-representation file more than one variable for each type,
- together with the other characteristics or prerequisites of entities (e.g., gender, or implied properties), so as to be able to generate more meanings of an ambiguous sentence.

# Concluding remarks

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- The problem of an anaphoric reference to a previously used expression is a well-known hard nut of *linguistic* analysis, because the antecedent of the anaphoric reference is often not unambiguously determined.
- Thus it is often said that anaphora constitutes a *pragmatic* problem rather than a problem of logical semantics.
- We agree that *logical* analysis cannot disambiguate any sentence, because it presupposes understanding and full linguistic competence.
- Yet our method of logical analysis *contributes to solving the problem of disambiguation* in at least two respects:
  - (a) the type-theoretical analysis often unambiguously determines which of the possible meanings of a homonymous expression is used in a sentence, and
  - (b) if there are two or more possible readings of a sentence, the logical analysis should make all of them explicit. This often concerns the distinction between *de dicto* and *de re* readings.
- In this sense *anaphora is a semantic problem*.

# Thank you for your attention

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- Questions, answers, comments, ...